

## On the ratio of two blocks of consecutive integers

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**Abstract.** Under certain assumptions, it is shown that eq. (2) has only finitely many solutions in integers  $x \geq 0$ ,  $y \geq 0$ ,  $k \geq 2$ ,  $l \geq 0$ . In particular, it is proved that (2) with  $a=b=1$ ,  $l=k$  implies that  $x=7$ ,  $y=0$ ,  $k=3$ .

**Keywords.** Exponential Diophantine equations; elementary Diophantine approximations; linear forms in logarithms; arithmetic-geometric mean.

### 1.

Let  $a$  and  $b$  be relatively prime positive integers. Erdős [5] conjectured that there are only finitely many integers  $x \geq 0$ ,  $y \geq 0$ ,  $k \geq 2$ ,  $l \geq 0$  with  $k+l \geq 3$  satisfying

$$x \geq y + l + k \tag{1}$$

and

$$a(x+1) \cdots (x+k) = b(y+1) \cdots (y+k+l). \tag{2}$$

By letters  $x, y, k$  and  $l$ , we shall always understand that  $x \geq 0$ ,  $y \geq 0$ ,  $k \geq 2$ ,  $l \geq 0$  with  $k+l \geq 3$  are integers satisfying (2). For an integer  $v$  with  $|v| > 1$ , we denote by  $P(v)$  and  $\omega(v)$ , respectively, the greatest prime factor of  $v$  and the number of distinct prime factors of  $v$ . Further, we write  $P(0)=P(1)=P(-1)=1$  and  $\omega(0)=\omega(1)=\omega(-1)=0$ .

In this and the next paragraph, we state some of the earlier results on (2). Mordell [7] proved that (2) with  $a=b=1$ ,  $k=2$ ,  $l=1$  implies that either  $x=1$ ,  $y=0$  or  $x=13$ ,  $y=4$ . Avanesov [1] confirmed a conjecture of Sierpinski by proving that  $x=0$ ,  $y=0$ ;  $x=3$ ,  $y=2$ ;  $x=14$ ,  $y=7$ ;  $x=54$ ,  $y=19$  and  $x=118$ ,  $y=33$  are the only solutions of (2) with  $a=3$ ,  $b=1$ ,  $k=2$ ,  $l=1$ . Tzanakis and de Weger [12] determined all the solutions of (2) with  $a=1$ ,  $b=2$ ,  $k=2$ ,  $l=1$ . Boyd and Kisilevsky [3] showed that  $x=1$ ,  $y=0$ ;  $x=3$ ,  $y=1$  and  $x=54$ ,  $y=18$  are the only solutions of (2) with  $a=b=1$ ,  $k=3$ ,  $l=1$ . Cohn [4] proved that (2) with  $a=1$ ,  $b=2$ ,  $k=4$ ,  $l=0$  is satisfied only if  $x=4$ ,  $y=3$ . Further, Ponnudurai [8] showed that  $x=2$ ,  $y=1$  and  $x=6$ ,  $y=4$  are the only solutions of (2) with  $a=1$ ,  $b=3$ ,  $k=4$ ,  $l=0$ . Shorey [10] showed that (2) with  $l=0$  and (1) implies that either  $k$  is bounded by an effectively computable number depending only on  $a$ ,  $b$  or  $k = [\alpha + 1]$  where

$$\alpha = \log \left( \frac{b}{a} \right) / \log \left( \frac{x}{y} \right).$$

Further, it is proved in [10] that (2) with  $l=0$  and (1) implies that  $k$  is bounded by an effectively computable number depending only on  $a, b, P(x), P(y)$  and also,

$$\log x \leq C_1 k \quad (3)$$

where  $C_1$  is an effectively computable number depending only on  $a, b, P(y)$  and  $P(x-y)$ . On the other hand, we see from Cramer's conjecture on distance between consecutive primes that (2) and (1) imply that

$$(\log x)^2 > C_2 k \quad (4)$$

where  $C_2 > 0$  is an absolute constant. In this paper, we shall extend these results to a more general equation (2). For a given  $k$  and  $l$ , we refer to a theorem of Siegel [11] to observe that (2) has only finitely many solutions in  $x$  and  $y$  provided that the curve represented by (2) is irreducible over the field of complex numbers and has positive genus.

It has not been possible to confirm the conjecture of Erdős, stated above, even when  $y$  is bounded. Erdős [5] considered a particular case of (2) corresponding to  $y=0$ , namely,

$$(X+1)\cdots(X+M) = N! \quad (5)$$

where  $X \geq 2, M \geq 2$  and  $N \geq 2$  are integers. Erdős[5] conjectured that  $8.9.10 = 6!$  is the only solution of (5). Erdős[5] proved that for  $\varepsilon > 0$  there exists  $N_0$  depending only on  $\varepsilon$  such that (5) with  $N \geq N_0$  implies that

$$X \geq (2-\varepsilon)^N. \quad (6)$$

We re-write (5) with  $M=2$  as

$$(2X+3)^2 - 1 = 4N!$$

which reminds us of the open problem on squares of the form  $N! + 1$ .

By fixing any three of the four variables  $x, y, k$  and  $l$  in (2), the fourth one is determined uniquely, if it exists. For given  $x, y$  and  $l$ , we start with the following result that determines the exact value of  $k$  if it exists.

**Theorem 1.** *Let  $y' = \max(y, 1)$  and*

$$\beta = \left( \log \left( \frac{b}{a} \right) + l \log y' \right) / \log \left( \frac{x}{y} \right). \quad (7)$$

*There exists an effectively computable number  $C_3$  depending only on  $a$  and  $b$  such that (2) with (1),  $k \geq C_3$  and*

$$y > (k+l)^2 \text{ if } 12 \log(l+1) \geq k \quad (8)$$

*implies that*

$$0 < k - \beta < 1. \quad (9)$$

We observe that a restriction of the type (8) is necessary for obtaining (9). For

$0 < \phi < (\log 2)/12$ ,  $l = [e^{\phi k}]$ ,  $0 < y \leq (k+l)^2$  and  $k$  exceeding a sufficiently large number depending only on  $a$ ,  $b$  and  $\phi$ , we see from (7), (15) and (23) that  $\beta < C_4 k$  where  $0 < C_4 < 1$  is a number depending only on  $a$ ,  $b$  and  $\phi$ .

Next, we turn to a more general situation than considered in Theorem 1. For given  $x$  and  $y$ , we denote by  $N(x, y)$  the number of pairs  $(k, l)$  satisfying (2) and (1). For given  $x$ ,  $y$  and  $l$ , there is at most one  $k$  satisfying (2). Therefore, we see from Lemma 1 that

$$N(x, y) \leq C_5 \log x \tag{10}$$

where  $C_5$  is an effectively computable number depending only on  $a$  and  $b$ . In the following result, we sharpen (10) whenever  $y$  is somewhat smaller than  $x$ .

**Theorem 2.(a)** *There exist effectively computable numbers  $C_6$  and  $C_7$  depending only on  $a$  and  $b$  such that for every  $x$  and  $y$  with*

$$\log y \leq C_6 \frac{\log x}{\log \log x}, \tag{11}$$

we have

$$N(x, y) \leq C_7. \tag{12}$$

(b) *Let  $\varepsilon > 0$  and  $y < (1 - \varepsilon)x$ . Then*

$$N(x, y) \leq C_8 \log \log x$$

where  $C_8$  is an effectively computable number depending only on  $a$ ,  $b$  and  $\varepsilon$ .

A pair  $(k, l)$  in Theorem 2 may depend on  $x$  and  $y$ . Now, for  $\varepsilon > 0$  and  $y \leq x^{1-\varepsilon}$ , we show that  $\max(k, l)$  is bounded by a number depending only on  $\varepsilon$ ,  $a$ ,  $b$  and  $P(x)$ . See Theorem 3(b) which finds an application in the proof of Theorem 5. Furthermore, we give lower bounds for  $P(x)$  and  $P(x-y)$  whenever  $y$  is smaller than some power of  $x$  and we apply these estimates in the proof of Corollary 1.

**Theorem 3.** *Suppose that (2) with (1) is satisfied. Let  $\varepsilon > 0$ .*

(a) *If*

$$x - y \geq \min \left( \varepsilon x, \frac{x((\log k)(\log \log k))^2}{k} \right), \tag{13}$$

then  $l$  is bounded by an effectively computable number depending only on  $\varepsilon$ ,  $a$ ,  $b$  and  $P(x)$ .

(b) *If  $y \leq x^{1-\varepsilon}$ , then  $\max(k, l)$  is bounded by an effectively computable number depending only on  $\varepsilon$ ,  $a$ ,  $b$  and  $P(x)$ .*

(c) *There exist effectively computable numbers  $\delta > 0$  and  $C_9$  depending only on  $a$  and  $b$  such that the inequalities*

$$y \leq x^\delta \quad \text{and} \quad P(x) \leq l + k$$

imply that  $\max(x, y, k, l) \leq C_9$ .

(d) *There exist effectively computable numbers  $\delta_1 > 0$  and  $C_{10}$  depending only on  $a$ ,  $b$  and  $P(x-y)$  such that if  $y \leq x^{\delta_1}$ , then  $\max(x, y, k, l) \leq C_{10}$ .*

(e) *There is  $C_{11}$  depending only on  $a$ ,  $b$  and  $y$  such that  $P(x-y) \leq l$  implies that  $\max(x, k, l) \leq C_{11}$ .*

The proof of Theorem 3(a) depends on the theory of linear forms in logarithms. We combine Theorem 3(c), (d) with Lemma 8. We derive that (2) with (1) and  $P(x) \leq l+k$  implies that  $l/k$  is bounded by an effectively computable number depending only on  $a$  and  $b$ . Also, we see that (2) and (1) imply that  $l/k$  is bounded by an effectively computable number depending only on  $a$ ,  $b$  and  $P(x-y)$ . Furthermore, if  $k$  is fixed, we may apply the theorem of Siegel, stated above, to finitely many pairs  $(k, l)$  given by the above two assertions to derive that (2) has only finitely many solutions in  $x$  and  $y$  under certain assumptions already mentioned. Next, we consider (2) with  $l/k = 1$  and  $a = b = 1$  i.e.

$$(x+1)\cdots(x+k) = (y+1)\cdots(y+2k). \quad (14)$$

MacLeod and Barrodale [6] showed that (14) with  $k \in \{2, 4, 5\}$  has no solution in  $x, y$  and (14) with  $k=3$  implies that  $x=7, y=0$ . Further, MacLeod and Barrodale [6] proved that for a given  $k$ , there are only finitely many pairs (depending on  $k$ )  $x, y$  satisfying (14). We prove

**Theorem 4.** *The equation (14) has only one solution in integers  $x \geq 0, y \geq 0, k \geq 2$  and it is given by  $x=7, y=0, k=3$ .*

In Theorems 2, 3 and 4, we see that  $y$  is somewhat smaller than  $x$ . In the next result, we replace this by the assumption that  $P(y)$  is bounded. We combine Theorem 3(b) with the theory of linear forms in logarithms to obtain the following result.

**Theorem 5.** *Equation (2) with (1) implies that  $\max(x, y, k, l)$  is bounded by an effectively computable number depending only on  $a, b, P(x)$  and  $P(y)$ .*

We may derive from Theorem 3(c) that (5) implies  $P(X) > N$  whenever  $X$  exceeds a sufficiently large effectively computable absolute constant. In fact, we apply Theorem 3(e) to obtain a more general result on (5).

#### COROLLARY 1.

*Let  $B \geq 0$  be an integer. There exists an effectively computable number  $C_{12}$  depending only on  $B$  such that (5) with  $N > C_{12}$  implies that  $P(X-B) > N$ .*

For the proofs of our Theorems, we prove certain estimates that are of independent interest. For example, we show unconditionally that (4) is valid whenever  $y < (1-\varepsilon)x$  for  $\varepsilon > 0$ . We formulate these estimates as the following theorem.

**Theorem 6.** *Suppose that (2) with (1) is satisfied. Then*

(a) *There exists an effectively computable number  $C_{13} > 0$  depending only on  $a$  and  $b$  such that*

$$x \geq C_{13} k^3 (\log k)^{-4}. \quad (15)$$

(b) *Let  $\varepsilon > 0$  and  $y < (1-\varepsilon)x$ . Then*

$$(\log x)^2 > C_{14} k \quad (16)$$

*where  $C_{14} > 0$  is an effectively computable number depending only on  $\varepsilon, a$  and  $b$ .*

(c) There exist effectively computable numbers  $C_{15}$  and  $C_{16} > 0$  depending only on  $a$  and  $b$  such that if  $k \geq C_{15}$  and  $y \leq (k+l)^3 (\log(k+l))^{-5}$ , we have

$$\log \log x > C_{16}k.$$

(d) There exists an effectively computable number  $C_{17} > 0$  depending only on  $a$  and  $b$  such that

$$x - y \geq C_{17}x^{2/3}.$$

## 2.

This section contains a proof of Theorem 6 and lemmas for our Theorems 1, 2, 3, 5. Throughout this section, we suppose that (1) and (2) are satisfied and we shall use this assumption without reference. We put

$$U_i = ax^i - by^{i+1} \quad \text{for } 0 \leq i \leq k \quad (17)$$

and

$$f(x, y) = x - (b/a)^{1/k} y^{1+(l/k)}. \quad (18)$$

Let

$$F(z) = (z+1) \cdots (z+k) = z^k + A_1 z^{k-1} + \cdots + A_k \quad (19)$$

and

$$G(z) = (z+1) \cdots (z+l) = z^l + A'_1 z^{l-1} + \cdots + A'_l. \quad (20)$$

Then, we refer to [6, p. 256] to observe that

$$0 < A_i \leq (k+1)^{2i}/2^i! \quad \text{for } 1 \leq i \leq k \quad (21)$$

and

$$0 < A'_j \leq (l+1)^{2j}/2^j! \quad \text{for } 1 \leq j \leq l. \quad (22)$$

We start with the following result that provides an upper bound for  $l$ .

*Lemma 1.* There exists an effectively computable absolute constant  $c$  such that

$$l \leq c \log(a+1) + (2 \log x)/\log 2. \quad (23)$$

*Proof.* We re-write (2) as

$$a \frac{(x+1) \cdots (x+k)}{k!} = b \frac{(y+1) \cdots (y+k+l)(k+l)!}{(k+l)! k!}. \quad (24)$$

We observe that

$$\text{ord}_2(\text{R.H.S. of (24)}) \geq (l-1)/2. \quad (25)$$

On the other hand,

$$\text{ord}_2(\text{L.H.S. of (24)}) \leq \text{ord}_2(a) + \max_{1 \leq i \leq k} \text{ord}_2(x+i) \leq (\log(ax+k))/\log 2. \quad (26)$$

Finally, we combine (24), (25) and (26) to derive (23). □

We give estimates for  $f(x, y)$  in the next three lemmas.

*Lemma 2.*

$$f(x, y) > 0. \tag{27}$$

*Proof.* We suppose that  $f(x, y) \leq 0$ . Then, we see from (17) and (18) that  $U_k \leq 0$ , which, since  $x > y$ , implies that

$$U_i \leq 0 \quad \text{for } 0 \leq i \leq k. \tag{28}$$

Now, we derive from (2), (19), (20), (17) and (28) that

$$0 < aF(x) - bF(y)y^l = U_k + A_1 U_{k-1} + \dots + A_k U_0 \leq 0 \tag{29}$$

which is a contradiction. □

By (2), we observe that  $ax^k < b(y+k+l)^{k+l}$  which implies that

$$x < (b/a)^{1/k} (y+k+l)^{1+(l/k)} \leq \theta (y+k+l)^{1+(l/k)} \tag{30}$$

where

$$\theta = \max(1, (b/a)^{1/k}) \leq \max(1, (b/a)).$$

Now, we give an upper estimate for  $f(x, y)$ .

*Lemma 3.* For  $l \leq k$ , we have

$$f(x, y) \leq 16\theta k (\max(k, y))^{l/k}.$$

*Proof.* Suppose that  $y > k+l$ . Then, by (30) and  $l \leq k$ , we have

$$\begin{aligned} x &< (b/a)^{1/k} y^{1+(l/k)} \left(1 + \frac{k+l}{y}\right)^{1+(l/k)} \\ &\leq (b/a)^{1/k} y^{1+(l/k)} \left(1 + \frac{(k+l)^2}{ky} + \frac{l(k+l)^3}{2k^2y^2}\right) \end{aligned}$$

which implies that

$$f(x, y) \leq 6\theta ky^{l/k}.$$

If  $y = k+l$ , we see from (30) and  $l \leq k$  that

$$f(x, y) \leq 3\theta y^{1+(l/k)} \leq 6\theta ky^{l/k}.$$

Thus, we may suppose that  $y < k+l$ . Then

$$\begin{aligned} x &< \theta(k+l)^{1+(l/k)} \left(1 + \frac{y}{k+l}\right)^{1+(l/k)} \\ &\leq \theta(k+l)^{1+(l/k)} \left(1 + \frac{y}{k} + \frac{ly^2}{2k^2(l+k)}\right) \end{aligned}$$

which implies that

$$f(x, y) < x \leq 16\theta k^{1+(l/k)}. \quad \square$$

For applications, it is convenient to formulate a version of Lemma 3 which is valid also for  $l > k$ .

*Lemma 4.*

$$f(x, y) \leq 16\theta kx^{l/k}. \quad (31)$$

*Proof.* By (1), we observe that  $x > \max(k, y)$ . Now, we apply Lemma 3 to assume that  $l > k$ . Then, the trivial estimate  $f(x, y) < x$  implies (31).  $\square$

We apply our estimates on  $f(x, y)$  to give bounds for  $x - y$ .

*Lemma 5.* There exist effectively computable numbers  $c_1$  and  $c_2 > 0$  depending only on  $a$  and  $b$  such that

$$x - y \leq \begin{cases} c_1 \left( k + \frac{lx \log x}{k} \right) & \text{if } l > 0 \\ c_1 \left( k + \frac{x}{k} \right) & \text{if } l = 0 \end{cases} \quad (32)$$

and

$$x - y \geq \begin{cases} (c_2 ly \log(y+1))/k, & \text{if } l > 0 \\ c_2 y/k & \text{if } l = 0. \end{cases} \quad (33)$$

*Proof.* We write  $c_3, c_4$  and  $c_5$  for effectively computable numbers depending only on  $a$  and  $b$ . We write, by (18),

$$x - y = f(x, y) + \Delta \quad (34)$$

where

$$\Delta = (b/a)^{1/k} y^{1+(l/k)} - y. \quad (35)$$

First, we prove (32). If  $l = 0$ , the assertion follows from (34), (35) with  $b > a$  and (31). Let  $l > 0$ . Then, we may assume that

$$l \log x < k, \quad (36)$$

otherwise (32) follows immediately. Now, we derive from (31) and (36) that

$$f(x, y) \leq 16\theta ek. \quad (37)$$

Further, it is easy to see from (35) and (36) that

$$\Delta \leq (c_3 lx \log x)/k. \quad (38)$$

Finally, we combine (34), (37) and (38) to obtain (32).

Next, we turn to the proof of (33). By (34) and (27), we observe that

$$x - y \geq \Delta. \quad (39)$$

If  $l=0$ , the assertion follows immediately from (39), (35) and  $b > a$ . Let  $l > 0$ . Then, we may suppose that  $y > 0$  and  $l \log(y+1) > c_4$  with  $c_4$  sufficiently large. Then, it is easy to see that  $\Delta \geq (c_3 l y \log y)/k$  which, together with (39), implies (33).

We apply an argument of Erodös and Lemmas 1, 5 to obtain a lower bound for  $x$  in the next lemma. The case  $l=0$  of this lemma is proved in [10].

*Lemma 6.* *There exists an effectively computable number  $c_6 > 0$  depending only on  $a$  and  $b$  such that*

$$x(\log x)^2 \geq \begin{cases} c_6 k^3 l^{-2} & \text{if } l > 0 \\ c_6 k^3 (\log k)^2 & \text{if } l = 0. \end{cases} \tag{46}$$

*Proof.* We denote by  $c_7, \dots, c_{13}$  effectively computable positive numbers depending only on  $a$  and  $b$ . We may assume that  $k \geq c_7$  with  $c_7$  sufficiently large and

$$x < k^3 \tag{41}$$

which, together with (23), implies that

$$l \leq c_3 \log k. \tag{42}$$

Furthermore, we derive from (2) and (1) that none of  $x+1, \dots, x+k$  is a prime number. Therefore, it follows from the well-known results on difference between consecutive primes that

$$x \geq k^{3/2}. \tag{43}$$

We denote by  $d$  the greatest common divisor of  $(x+1) \cdots (x+k)$  and  $(y+1) \cdots (y+k+l)$ . Then, by (2), we see that

$$x^k < (x+1) \cdots (x+k) \leq bd. \tag{44}$$

Let  $S = \{x+1, \dots, x+k\}$ . For a prime  $p \leq k$ , we choose an  $f(p) \in S$  such that  $p$  does not appear to a higher power in the factorization of any other element of  $S$ . Let  $S_1$  be the subset of  $S$  obtained by deleting all  $f(p)$  with  $p \leq k$ . Then, by a fundamental argument of Erdös, we have

$$\prod_{s \in S_1} s \leq \prod_{p \leq k} p^{\lceil (k/p) \rceil + \lceil (k/p^2) \rceil + \dots} = k!.$$

Therefore, the contribution  $d_1$  in  $d$  from all primes not exceeding  $k$  is at most

$$k^k (x+k)^{\pi(k)} \leq (e^k k)^k,$$

by (41). Further, the contribution  $d_2$  in  $d$  from all primes  $p$  with  $k < p \leq (2k+l-1)$  is less than or equal to

$$(x+k)^{\pi(2k+l-1) - \pi(k)} \leq e^{4k},$$

by (41) and (42). Now, we set

$$d_3 = d/d_1 d_2 \text{ and } \Delta_1 = \left( \prod_{\mu = -(k+l-1)}^{k-1} (x-y+\mu) \right) / (2k+l-1)!.$$



Notice that  $\Delta_1$  is a positive integer not exceeding

$$e^{3k} \left( \frac{x-y}{2k} + 1 \right)^{2k+l-1},$$

by (42). Also, we observe from (2) that  $d_3 | \Delta_1$ . Consequently,

$$d \leq (c_9 k)^k \left( \frac{x-y}{2k} + 1 \right)^{2k+l-1}$$

which, together with (44) and (42), implies that

$$x \leq c_{10} k \left( \frac{x-y}{2k} + 1 \right)^{2+c_8(\log k)/k} \tag{45}$$

If  $l > 0$ , we combine (45), (43) and (32) to conclude that

$$x \leq c_{11} k \left( \frac{l x \log x}{k^2} \right)^{2+c_8(\log k)/k} \leq c_{12} k \left( \frac{l x \log x}{k^2} \right)^2,$$

by (41) and (42). If  $l = 0$ , we obtain in a similar way that  $x \leq c_{13} k (x/k^2)^2$ . The preceding two inequalities imply (40) immediately.  $\square$

Now, we are ready to prove Theorem 6(a), (b), (d).

*Proof of Theorem 6(a), (b), (d).* First, we observe that (15) is an immediate consequence of (40) and (23). We re-write  $y < (1-\varepsilon)x$  as  $x-y > \varepsilon x$  which, together with (32), (15) and (23), implies (16). Finally, we apply (33) and (40) to obtain Theorem 6(d).  $\square$

For applications, it is convenient to combine (32) and (15) to formulate the following result.

**COROLLARY 2.**

There exists an effectively computable number  $c_{14}$  depending only on  $a$  and  $b$  such that

$$x-y \leq (c_{14}(l+1)x \log x)/k. \tag{46}$$

For the proof of Theorem 6(c), we apply (15), (16) and (23) to obtain the following result which also finds an application in the proof of Theorem 1.

*Lemma 7.* Let  $\chi \geq 1$  and  $\gamma > 4/\log 2$ . There exist effectively computable numbers  $c_{15}$  and  $c_{16}$  depending only on  $a, b, \chi$  and  $\gamma$  such that for  $k \geq c_{15}$  and

$$\log(l+1) < (\gamma\chi)^{-1}k, \tag{47}$$

we have

$$y \geq \begin{cases} c_{16} l^x & \text{if } l > k \\ c_{16} k^3 (\log k)^{-4} & \text{if } l \leq k. \end{cases}$$

*Proof.* We may assume that  $c_{15}$  is sufficiently large. Let  $l > k$ . Then, we may suppose that  $y \leq l^x$ . Now, we derive from (30) and (23) that

$$\log x \leq c_{17} + (2l \log(3l^k))/k$$

where  $c_{17}$  and the subsequent letter  $c_{18} > 0$  are effectively computable numbers depending only on  $a$  and  $b$ . Therefore, by (23), we derive that  $k \leq (\gamma\chi) \log l$  which contradicts (47). Thus, we may suppose that  $l \leq k$ . Now, we may assume that  $y \leq k^3 (\log k)^{-4}$ . Then, we observe from (30) that  $k \geq x^{1/6}$  which, by Theorem 6(b) with  $\varepsilon = 1/2$  and (15), implies that

$$y \geq x/2 \geq c_{18} k^3 / (\log k)^4. \quad \square$$

*Proof of Theorem 6(c).* We may suppose that  $C_{15}$  is sufficiently large. Then, we apply Lemma 7 with  $\chi = 3$ ,  $\gamma = 6$  and (23) to obtain

$$c_{19} \log \log x \geq \log(l + 1) \geq k/18. \quad \square$$

Next, we prove a lemma which tells that certain assumptions involving variables of (2) are equivalent.

*Lemma 8(a).* Let  $0 < \delta < 1$ . There exist effectively computable numbers  $v_1$  and  $v_2 > 0$  depending only on  $\delta$  such that  $l \geq v_2 k$  whenever  $x \geq v_1$  and  $y \leq x^\delta$ .

(b) Let  $\mu > 0$ . There exist effectively computable numbers  $v_3$  and  $v_4$  with  $0 < v_4 < 1$  depending only on  $\mu$  such that  $y \leq x^{v_4}$  whenever  $x \geq v_3$  and  $l \geq \mu k$ .

*Proof.* (a). We may assume that  $v_1$  is sufficiently large. Let  $y \leq x^\delta < x/2$ . Then, we apply Theorem 6(b) with  $\varepsilon = 1/2$  to obtain (16) which, together with (30) and (23), implies that  $l \geq v_2 k$ .

(b) By (2),  $y < (a/b)^{1/(k+l)} (x+k)^{k/(k+l)}$  which implies the assertion. □

Finally, we state an estimate of Baker [2] on linear forms in logarithms and its  $p$ -adic analogue, due to Yu [13]. Let  $\alpha_1, \dots, \alpha_n$  be non-zero rational numbers of heights\* not exceeding  $A_1, \dots, A_n$ , respectively. We assume that  $A_j \geq 3$  for  $1 \leq j \leq n$ . We put

$$\Omega = \prod_{j=1}^n \log A_j, \quad \Omega' = \Omega / \log A_n.$$

Then we have

*Lemma 9.* (Baker [2]). There exists an effectively computable number  $c_{20}$  depending only on  $n$  such that the inequalities

$$0 < |\alpha_1^{b_1} \dots \alpha_n^{b_n} - 1| < \exp(-c_{20} \Omega \log \Omega' \log B)$$

have no solution in rational integers  $b_1, \dots, b_n$  of absolute values not exceeding  $B (\geq 2)$ .

*Lemma 10.* (Yu [13]). Let  $p$  be a prime number. Suppose that  $b_1, \dots, b_{n-1}$  and  $b_n = -1$  are rational integers of absolute values not exceeding  $B (\geq 2)$ . There exists an effectively computable number  $c_{21}$  depending only on  $n$  and  $p$  such that either  $\alpha_1^{b_1} \dots \alpha_n^{b_n} = 1$  or

$$\text{ord}_p(\alpha_1^{b_1} \dots \alpha_n^{b_n} - 1) \leq c_{21} \Omega \log \Omega' \log B.$$

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\*The height of a rational number  $u_1/u_2$  with  $\text{gcd}(u_1, u_2) = 1$  is defined as  $\max(|u_1|, |u_2|)$ .

3.

*Proof of Theorem 1.* We denote by  $c_{22}, c_{23}, \dots, c_{29}$  effectively computable positive numbers depending only on  $a$  and  $b$ . We may assume that  $k \geq c_{22}$  with  $c_{22}$  sufficiently large. Let  $F(z)$ ,  $G(z)$  and  $U_i$  with  $0 \leq i \leq k$  be given by (19), (20) and (17), respectively. We apply Lemma 7 with  $\chi > 2$ ,  $4/\log 2 < \gamma < 6$  and  $\gamma\chi = 12$  to derive from (8) that

$$y > (k+l)^2. \tag{48}$$

By (27), we observe that  $U_k > 0$ . Therefore, it suffices to show that  $U_{k-1} < 0$ . We assume that

$$U_{k-1} \geq 0 \tag{49}$$

and we shall arrive at a contradiction. By (2),

$$aF(x) - bF(y)y^l = bF(y)R(y) \tag{50}$$

where  $R(y) = 0$  if  $l = 0$  and for  $l > 0$ ,

$$R(y) = G(y+k) - y^l = R_1(y) + R_2(y)$$

with

$$R_1(y) = (y+k)^l - y^l, \quad R_2(y) = A'_1(y+k)^{l-1} + \dots + A'_l.$$

Then, it is easy to derive from the estimates (21), (22) and (48) that

$$F(y) \leq 3y^k, \quad R_1(y) \leq 2kly^{l-1}, \quad R_2(y) \leq 2l^2y^{l-1}.$$

Therefore, we derive from (50) that

$$aF(x) - bF(y)y^l = bF(y)R(y) \leq 6bl(k+l)y^{k+l-1}. \tag{51}$$

We notice from (17) that

$$U_k - xU_{k-1} = by^{k+l-1}(x-y) > 0$$

which, by (49), (29) and (50), implies that

$$by^{k+l-1}(x-y) \leq U_k = bF(y)R(y) - A_1U_{k-1} - \dots - A_kU_0. \tag{52}$$

Now, by (49), (21) and (48), we derive that

$$-A_1U_{k-1} - \dots - A_kU_0 \leq -A_2U_{k-2} - \dots - A_kU_0 \leq c_{23}k^4y^{k+l-2}$$

which, together with (51) and (52), implies that

$$x - y \leq c_{24}(kl + l^2 + k^4y^{-1}). \tag{53}$$

If  $l = 0$ , we see from (33) and (53) that  $y \leq k^{11/4}$  which contradicts Lemma 7. Thus, we may assume that  $l > 0$ . Now, we combine again (33) and (53) to derive that

$$y \log y \leq c_{25}(k^2 + kl + k^5(ly)^{-1}). \tag{54}$$

Now, we combine (54) and (48) to observe that  $l^2 \leq c_{25}(k^2 + kl + k^3l^{-1})$  which implies that

$$l \leq c_{26}k. \tag{55}$$

Now, we see from (54), (48) and (55) that

$$y^2 \log y \leq c_{27}k^5l^{-1}. \tag{56}$$

Then, we apply Theorem 6 to derive that

$$(\log x)^2 \geq c_{28}k. \tag{57}$$

Finally, we combine (1), (30), (57), (56) and (55) to conclude that  $k < x \leq c_{29}$ .  $\square$

*Proof of Theorem 2 (a).* We write  $c_{30}, c_{31}, \dots, c_{37}$  for effectively computable positive numbers depending only on  $a$  and  $b$ . Suppose that the assertion (12) is not valid. For a given  $k$ , we observe that there is at most one  $l$  satisfying (2). Therefore, we may assume that (1) and (2) are satisfied with  $k=k_1, l=l_1$  and  $k=k_2, l=l_2$  such that  $k_1 < k_2, k_1 \geq c_{30}$  and  $k_2 - k_1 \geq c_{30}$  with  $c_{30}$  sufficiently large. Then, by (2) with  $k=k_1, l=l_1$  and  $k=k_2, l=l_2$ , we observe that

$$(X + 1) \cdots (X + k_2 - k_1) = (Y + 1) \cdots (Y + k_2 - k_1 + l_2 - l_1) \tag{58}$$

where

$$X = x + k_1, \quad Y = y + k_1 + l_1. \tag{59}$$

By (59) and (1), we see that  $X > Y$  which, together with (58), implies that  $l_1 < l_2$ . Further, by (30), (11), (23) and Theorem 6 (b), we derive that

$$k_i \log \log x \leq c_{31}(l_i + 1) \quad (i = 1, 2) \tag{60}$$

which, together with (23), implies that

$$k_i \leq c_{32}(\log x) / \log \log x \quad (i = 1, 2). \tag{61}$$

Now, we see from (58), (59), (11), (61) and (23) that

$$(k_2 - k_1) \log \log x \leq c_{33}(l_2 - l_1). \tag{62}$$

By counting the power of 2 on both the sides of (24) with  $a=b=1, x=X, y=Y, k=k_2 - k_1$  and  $l=l_2 - l_1$  obtained from (58), we have

$$l_2 - l_1 \leq 2 \max_{1 \leq i \leq k_2 - k_1} \text{ord}_2(X + i) + 1. \tag{63}$$

We show that

$$\text{either } l_1 \leq c_{34} \log k_2 \quad \text{or} \quad l_2 - l_1 \leq c_{34} \log k_2. \tag{64}$$

If

$$\max_{1 \leq j \leq k_1} \text{ord}_2(x + j) = \text{ord}_2(x + j_0) \leq 2 \log k_2, \tag{65}$$

then we count the power of 2 on both the sides in (24) with  $k=k_1, l=l_1$  to obtain

$l_1 \leq c_{35} \log k_2$ . Therefore, we may suppose that (65) is not valid. Then, by (59), we write

$$X + i = x + j_0 + k_1 + i - j_0$$

to observe that

$$\text{ord}_2(X + i) = \text{ord}_2(k_1 + i - j_0) \leq 2 \log k_2 \tag{66}$$

for  $1 \leq i \leq k_2 - k_1$ . Then, we see from (63) and (66) that  $l_2 - l_1 \leq 5 \log k_2$ . This proves (64). Now, we combine (60), (62), (64) and (61) to derive that either  $k_1 \leq c_{36}$  or  $k_2 - k_1 \leq c_{36}$  which is not possible if  $c_{30} > c_{36}$ .

(b) Suppose that (1) and (2) with  $k = k_1, l = l_1$  and  $k = k_2, l = l_2$  are satisfied. Then, we observe that (64) is valid. Now, we apply Theorem 6(b) to derive that either  $l_1 \leq c_{37} \log \log x$  or  $l_2 - l_1 \leq c_{37} \log \log x$  where  $c_{37}$  is an effectively computable number depending only on  $\varepsilon, a$  and  $b$ . Finally, we observe that for a given  $l$  there is at most one  $k$  satisfying (2) to complete the proof of Theorem 2(b).  $\square$

4.

*Proof of Theorem 3(a).* We denote by  $d_1, \dots, d_5$  effectively computable positive numbers depending only on  $\varepsilon, a, b$  and  $P(x)$ . We may assume that  $x > d_1$  with  $d_1$  sufficiently large, otherwise the assertion follows from (1). By (24), (25) and (26), we see that

$$l \leq d_2 \max_{1 \leq i \leq k} \text{ord}_2(x + i) \leq d_2 \left( \max_{1 \leq i \leq k} \text{ord}_2(-xi^{-1} - 1) + \frac{\log k}{\log 2} \right).$$

Further, we write  $x$  as  $\prod_{j=1}^{\omega(x)} p_j^{\text{ord}_{p_j}(x)}$  where  $p_1, \dots, p_{\omega(x)}$  are the distinct prime factors of  $x$ . Now, we apply Lemma 10 to  $\text{ord}_2(-xi^{-1} - 1)$  with  $p = 2, n = \omega(x) + 2 \leq P(x) + 2, \alpha_1 = -1, b_1 = 1; \alpha_{j+1} = p_j, b_{j+1} = \text{ord}_{p_j}(x)$  with  $1 \leq j \leq \omega(x)$  and  $\alpha_n = i, b_n = -1$ , to obtain

$$l \leq d_3 (\log \log x) \log k$$

which, together with (46) and (13), implies that

$$(\log k)(\log \log k)^2 \leq d_4 (\log x)(\log \log x).$$

Therefore, since  $\omega(x) \leq P(x)$ , we have

$$x > k^{\omega(x)}. \tag{67}$$

Consequently, there exists a prime  $p$  dividing  $x$  such that

$$p^{\text{ord}_p(x)} > k. \tag{68}$$

Now, we count the power of  $p$  on both the sides of (24) to derive that

$$[l/p] - 1 \leq \text{ord}_p(a) \tag{69}$$

which implies that  $l \leq d_5$ .  $\square$

(b) We apply Lemma 8(a) with  $\delta = 1 - \varepsilon$  to conclude that  $k \leq d_6 l$  where  $d_6$  and the subsequent letter  $d_7$  are effectively computable numbers depending only on  $\varepsilon$ ,  $a$ ,  $b$  and  $P(x)$ . Now, we apply Theorem 3(a) to conclude that  $\max(k, l) \leq d_7$ .  $\square$

(c) We write  $d_8, \dots, d_{13}$  for effectively computable numbers depending only on  $a$  and  $b$ . There is no loss of generality in assuming that  $x > d_8$  with  $d_8$  sufficiently large. Suppose that (2) with  $P(x) \leq l + k$  and  $y \leq x/2$  is satisfied. For a prime  $p$  with  $k < p \leq k + l$  and  $p|x$ , we derive from (2) that  $p|a$ . Therefore

$$P(x) \leq \max(k, P(a)).$$

Consequently, we observe from prime number theory that

$$\omega(x) \leq d_9 k / \log k. \quad (70)$$

First, we show that

$$l \leq d_{10} k. \quad (71)$$

Let  $x \leq k^{\omega(x)}$ . Then, we see from (23) and (70) that  $l \leq d_{11} k$ . Thus, we may assume (67) which implies (68), (69) and hence,  $l \leq d_{12} k$ . This proves (71). Finally, we combine (30), (71) and Theorem 6(b) with  $\varepsilon = 1/2$  to conclude that  $y > x^{d_{13}}$ .

(d) Let  $d_{14}, \dots, d_{18}$  be effectively computable positive numbers depending only on  $a$ ,  $b$  and  $P(x - y)$ . We may assume that  $x > d_{14}$  with  $d_{14}$  sufficiently large. Suppose that (2) with  $y < x/2$  is satisfied. Then, the inequality (16) is valid. Further, we may suppose that

$$l \geq \max(P(x - y), 2a), \quad (72)$$

otherwise, we may derive from (30) and (16) that  $y \geq x^{d_{15}}$ .

Let  $F(z)$  be given by (19). We re-write (2) as

$$0 \neq aF(x) - aF(y) = F(y)(b(y + k + 1) \cdots (y + k + l) - a). \quad (73)$$

For a prime  $p$  dividing  $x - y$ , we see from (73) and (72) that

$$\text{ord}_p(x - y) \leq \text{ord}_p(F(y)) + \text{ord}_p(a). \quad (74)$$

Further, we observe that

$$\text{ord}_p(F(y)) \leq \max_{1 \leq i \leq k} \text{ord}_p(y + i) + \left[ \frac{k}{p} \right] + \left[ \frac{k}{p^2} \right] + \dots \quad (75)$$

Now, we derive from  $y < x/2$ , (74) and (75) that

$$\log \left( \frac{x}{2} \right) \leq \log(x - y) = \sum_{p|(x-y)} \text{ord}_p(x - y) \log p \leq d_{16} (\log(y + k) + k)$$

which, by (16) and (23), implies that either  $y \geq x^{d_{17}}$  or  $l \leq d_{17} k$ . Finally, as above, we apply (30) and (16) to assume that the latter inequality is not valid.

(e) We may assume (72). Therefore, the inequality (74) is valid. Consequently, there

is an effectively computable number  $d_{18}$  depending only on  $a, b$  and  $y$  such that

$$\begin{aligned} \log\left(\frac{x}{2}\right) &\leq \log(x-y) \leq \log a + \pi(l) \log(y+k) + 2k \sum_{p \leq l} \frac{\log p}{p} \\ &\leq d_{18} \left( \frac{l \log k}{\log l} + k \log l \right) \end{aligned}$$

which, together with (23) and Theorem 6(c), completes the proof. □

*Proof of Theorem 5.* We denote by  $d_{19}, \dots, d_{28}$  effectively computable positive numbers depending only on  $a, b, P(x)$  and  $P(y)$ . We may assume that  $x > d_{19}$ , with  $d_{19}$  sufficiently large, otherwise the theorem follows from (1). We apply Lemma 9 to conclude that

$$x - y \geq x (\log x)^{-d_{20}}. \tag{76}$$

On the other hand, we derive from (46) and (23) that

$$x - y \leq d_{21} x (\log x)^2 / k. \tag{77}$$

We combine (76) and (77) to derive that

$$k \leq (\log x)^{d_{22}}. \tag{78}$$

Now, we show that

$$y \leq (\log x)^{d_{23}}. \tag{79}$$

For proving (79), we refer to (23) and (78) to assume that

$$y > (k + l)^4. \tag{80}$$

Then, by (27), (2), (15) and (80), we observe that

$$0 < U_k \leq d_{24} ((k + l)^2 y^{k+l-1} + k^2 x^{k-1}) \tag{81}$$

where  $U_k$  is given by (17). On the other hand, we apply again Lemma 9 to obtain

$$U_k \geq \max(x^k, y^{k+l}) ((k + l) \log x)^{-d_{25}}$$

which, together with (78) and (23), implies that

$$U_k \geq \max(x^k, y^{k+l}) (\log x)^{-d_{26}}. \tag{82}$$

Finally, we combine (81), (82), (78) and (80) to obtain (79). Therefore, we conclude from Theorem 3(b) with  $\varepsilon = 1/2$  to obtain that  $\max(k, l) \leq d_{27}$  which, together with (30) and (79), implies that  $x = \max(x, y, k, l) \leq d_{28}$ . □

*Proof of Corollary 1.* We denote by  $d_{29}, d_{30}$  and  $d_{31}$  effectively computable positive numbers depending only on  $B$ . We may assume that  $N \geq d_{29}$ , with  $d_{29}$  sufficiently large. In (2), we put  $a = 1, x = X, k = M, b = B!, y = B$  and  $l = N - M - B$ . By (6), we notice that  $X \geq N$  so that (1) is satisfied. Therefore, we derive from Theorem 6(c) that

$$M \leq d_{30} \log \log X. \tag{83}$$

Further, we see from (5) that

$$N \geq d_{3,1}(\log X)/\log \log X. \tag{84}$$

Now, we apply Theorem 3(e), (83) and (84) to conclude that

$$P = :P(X - B) > N - M - B > M + B. \tag{85}$$

Further, we see from (5) and (85) that

$$P \left| \left( \frac{N!}{M!} - \frac{(B+1)\cdots(B+M)}{M!} \right) \right|. \tag{86}$$

Finally, we derive from (86) and (85) that  $P > N$ . □

### 5.

This section is devoted to preliminaries for the proof of Theorem 4. Let  $F(z)$  be given by (19). Then

$$A_j = \sum_{i=i_0=1}^k \sum_{i_1=1}^{i_0-1} \cdots \sum_{i_{j-1}=1}^{i_{j-2}-1} i_0 i_1 \cdots i_{j-1} \quad \text{for } 1 \leq j \leq k.$$

We write the right hand side of (14)

$$(y+1)\cdots(y+2k) = \prod_{j=1}^k (u+j(2k-j+1)) \tag{87}$$

where

$$u = y(y+2k+1). \tag{88}$$

Let

$$G(z) = \prod_{j=1}^k (z+j(2k-j+1)) = z^k + B_1 z^{k-1} + \cdots + B_k \tag{89}$$

where

$$B_j = \sum_{i=i_0=1}^k \sum_{i_1=1}^{i_0-1} \cdots \sum_{i_{j-1}=1}^{i_{j-2}-1} i_0 i_1 \cdots i_{j-1} (2k-i_0+1)\cdots(2k-i_{j-1}+1) \tag{90}$$

for  $1 \leq j \leq k$ .

Further, we put

$$\Delta_q = \Delta_{q,k} = \sum_{i=1}^k i^q \quad \text{for } q = 1, 2, \dots$$

Then, we have  $A_1 = \Delta_1$ ,  $A_2 = (\Delta_3 - \Delta_2)/2$ ,

$$A_3 = (3\Delta_5 - 10\Delta_4 + 9\Delta_3 - 2\Delta_2)/24,$$

$$A_4 = (\Delta_7 - 7\Delta_6 + 17\Delta_5 - 17\Delta_4 + 6\Delta_3)/48,$$

$$B_1 = (2k+1)\Delta_1 - \Delta_2, \quad 3B_2 = \Delta_5 - (5k+4)\Delta_4 + (6k^2 + 12k + 5)\Delta_3 - (6k^2 + 7k + 2)\Delta_2$$



and

$$\begin{aligned}
 3B_3 = & -\frac{1}{6}\Delta_8 + \left(\frac{4}{3}k + \frac{22}{15}\right)\Delta_7 - \left(\frac{7}{2}k^2 + \frac{91}{10}k + \frac{149}{30}\right)\Delta_6 \\
 & + \left(3k^3 + \frac{35}{2}k^2 + \frac{137}{6}k + \frac{49}{6}\right)\Delta_5 - \left(10k^3 + \frac{59}{2}k^2 + \frac{51}{2}k + \frac{20}{3}\right)\Delta_4 \\
 & + \left(9k^3 + \frac{37}{2}k^2 + \frac{71}{6}k + \frac{71}{30}\right)\Delta_3 - \left(2k^3 + 3k^2 + \frac{7}{5}k + \frac{1}{5}\right)\Delta_2.
 \end{aligned}$$

Further, we derive from [9, p. 6] that

$$\begin{aligned}
 \Delta_1 &= k(k+1)/2, \quad \Delta_2 = k(k+1)(2k+1)/6, \quad \Delta_3 = k^2(k+1)^2/4, \\
 \Delta_4 &= k(k+1)(6k^3 + 9k^2 + k - 1)/30, \quad \Delta_5 = k^2(k+1)(2k^3 + 4k^2 + k - 1)/12, \\
 \Delta_6 &= k(k+1)(6k^5 + 15k^4 + 6k^3 - 6k^2 - k + 1)/42, \\
 \Delta_7 &= k^2(k+1)(3k^5 + 9k^4 + 5k^3 - 5k^2 - 2k + 2)/24
 \end{aligned}$$

and

$$\Delta_8 = k(k+1)(10k^7 + 35k^6 + 25k^5 - 25k^4 - 17k^3 + 17k^2 + 3k - 3)/90.$$

Consequently, we obtain the following expressions for  $A_1, A_2, A_3, A_4$  and  $B_1, B_2, B_3$ .

*Lemma 11.*  $A_1 = k(k+1)/2, A_2 = k(k+1)(k-1)(3k+2)/24,$

$$A_3 = k^2(k+1)^2(k-1)(k-2)/48,$$

$$A_4 = k(k+1)(k-1)(k-2)(k-3)(15k^3 + 15k^2 - 10k - 8)/5760 \tag{91}$$

and

$$B_1 = k(k+1)(2k+1)/3, \quad B_2 = k(k+1)(20k^4 + 16k^3 - 11k^2 - 19k - 6)/90, \tag{92}$$

$$\begin{aligned}
 B_3 = & k(k+1)(280k^7 - 28k^6 - 830k^5 - 745k^4 \\
 & + 136k^3 + 557k^2 + 486k + 144)/5670.
 \end{aligned}$$

Let

$$f = \left\lfloor \frac{2k^2}{3} \right\rfloor + k - \left\lfloor \frac{k}{2} \right\rfloor.$$

We write

$$F(z+f) = z^k + A_1(f)z^{k-1} + \dots + A_k(f) \tag{93}$$

where

$$A_i(f) = \binom{k}{i} f^i + \binom{k-1}{i-1} f^{i-1} A_1 + \dots + \binom{k-i+1}{1} f A_{i-1} + A_i \tag{94}$$

for  $1 \leq i \leq k$ .

Then, we apply Lemma 11 to obtain the following result.

*Lemma 12.* Let

$$\delta = \begin{cases} 0 & \text{if } k \equiv 0 \pmod{6} \\ -1/6 & \text{if } k \equiv 1, 5 \pmod{6} \\ -2/3 & \text{if } k \equiv 2, 4 \pmod{6} \\ 1/2 & \text{if } k \equiv 3 \pmod{6} \end{cases} \tag{95}$$

Then

$$A_1(f) - B_1 = \left(\frac{1}{6} + \delta\right)k,$$

$$A_2(f) - B_2 = \frac{2}{45}k^5 + \left(\frac{1}{9} + \frac{2\delta}{3}\right)k^4 + \left(-\frac{1}{24} + \frac{\delta}{3}\right)k^3 + \left(-\frac{7}{72} + \frac{\delta}{2}(\delta - 1)\right)k^2 \\ + \left(-\frac{1}{60} - \frac{\delta}{2}(1 + \delta)\right)k$$

and

$$A_3(f) - B_3 = \frac{4}{135}k^8 + \left(\frac{79}{2835} + \frac{2}{9}\delta\right)k^7 - \frac{5}{36}k^6 + \left(-\frac{169}{1080} - \frac{13}{18}\delta + \frac{1}{3}\delta^2\right)k^5 \\ + \left(\frac{67}{720} - \frac{17}{24}\delta - \frac{1}{2}\delta^2\right)k^4 + \left(\frac{1013}{6480} + \frac{9}{24}\delta - \frac{7}{12}\delta^2 + \frac{1}{6}\delta^3\right)k^3 \\ + \left(\frac{1}{72} + \frac{2}{3}\delta + \frac{1}{4}\delta^2 - \frac{1}{2}\delta^3\right)k^2 \\ + \left(-\frac{8}{315} + \frac{1}{6}\delta + \frac{1}{2}\delta^2 + \frac{1}{3}\delta^3\right)k.$$

*Proof.* We write  $g = \frac{2}{3}k^2 + \frac{1}{2}k$  and  $f - g = \delta$ . Then, we check that (95) is satisfied.

Further, we see from (94) that  $A_1(f) = kg + A_1 + \delta k$ ,  $A_2(f) = \binom{k}{2}g^2 + (k-1)A_1g + A_2 + \delta \binom{k}{2}(2g + \delta) + \delta(k-1)A_1$  and  $A_3(f) = \binom{k}{3}g^3 + \binom{k-1}{2}A_1g^2 + \binom{k-2}{1}A_2g + A_3 + \delta \binom{k}{3}(3g^2 + 3\delta g + \delta^2) + \delta \binom{k-1}{2}A_1(2g + \delta) + \delta \binom{k-2}{1}A_2$ . Now, we apply Lemma 11 to complete the proof of Lemma 12.  $\square$

Finally, we obtain estimates for all  $A_i(f)$  and  $B_j$ .

*Lemma 13.* For  $1 \leq i \leq k$ , we have

$$A_i(f) \leq (1 + \phi_k) \binom{k}{i} f^i \quad (96)$$

where

$$\phi_k = \frac{3}{4} \left(1 + \frac{1}{k}\right)^2 \left(1 - \frac{3}{8} \left(1 + \frac{1}{k}\right)^2\right)^{-1}. \quad (97)$$

Furthermore,

$$B_j \leq k^j(k+1)^{2j}/j! \quad \text{for } 1 \leq j \leq k \quad (98)$$

and

$$B_4 \leq (2k)^4 A_4, \quad B_4 \geq (k+1)^4 A_4. \quad (99)$$

*Proof.* By (94), we see that

$$A_i(f) \leq \binom{k}{i} f^i \left(1 + \frac{A_1}{f} \frac{i}{k} + \frac{A_2}{f^2} \frac{i(i-1)}{k(k-1)} + \cdots + \frac{A_i}{f^i} \frac{i!}{k(k-1)\cdots(k-i+1)}\right). \quad (100)$$

Further, by (21) and  $f > (2k^2)/3$ , we observe that

$$A_j f^{-j} \leq \frac{1}{j!} \left(\frac{3}{4}\right)^j \left(1 + \frac{1}{k}\right)^{2j}$$

which, together with (100), implies (96) where  $\phi_k$  is given by (97). By (90), we see that  $B_4 \geq (k+1)^4 A_4$  and

$$B_j \leq (2k)^j A_j, \quad \text{for } 1 \leq j \leq k,$$

which, by (21), implies (98). □

**6.**

In this section, we shall prove Theorem 4. The computations for the proof of Theorem 4 are carried out on a pocket calculator. Throughout this section, we assume that (14) is satisfied. For a prime  $p > 0$ , we write

$$n = b_0 + b_1 p + \dots + b_\mu p^\mu$$

where  $b_0, \dots, b_\mu$  are integers satisfying  $b_\mu \neq 0$  and  $0 \leq b_i < p$  for  $0 \leq i \leq \mu$ . Then, we start with the following well-known result.

*Lemma 14.*

$$\text{ord}_p(n!) = \frac{n - (b_0 + \dots + b_\mu)}{p - 1}. \tag{101}$$

For the proof of Lemma 14, we observe that

$$\text{ord}_p(n!) = \sum_{i=1}^{\mu} \left[ \frac{n}{p^i} \right] = \sum_{i=1}^{\mu} (b_i + b_{i+1} p + \dots + b_\mu p^{\mu-i})$$

which implies (101). We denote by  $\tau$  the number of positive integers  $i$  such that

$$\left[ \frac{y + 2k}{2^i} \right] - \left[ \frac{y}{2^i} \right] - \left[ \frac{2k}{2^i} \right] > 0.$$

Then, we apply Lemma 14 to obtain the following lower bound for  $x$  and this is fundamental for our argument.

*Lemma 15.*

$$x \geq 2^{k+\tau} - k. \tag{102}$$

Since  $\tau \geq 0$ , the inequality (102) implies that

$$x \geq 2^k - k. \tag{103}$$

*Proof.* We re-write (14) as

$$\frac{(x+1) \cdots (x+k)}{k!} = \frac{(y+1) \cdots (y+2k)(2k)!}{(2k)! k!}. \tag{104}$$

By (101),

$$\text{ord}_2\left(\frac{(2k)!}{k!}\right) = k. \tag{105}$$

Further, we observe that

$$\text{ord}_2\left(\frac{(y+1)\cdots(y+2k)}{(2k)!}\right) = \sum_{i=1}^{\infty} \left( \left\lfloor \frac{y+2k}{2^i} \right\rfloor - \left\lfloor \frac{y}{2^i} \right\rfloor - \left\lfloor \frac{2k}{2^i} \right\rfloor \right)$$

and every summand is either 0 or 1. Therefore

$$\text{ord}_2\left(\frac{(y+1)\cdots(y+2k)}{(2k)!}\right) = \tau. \tag{106}$$

On the other hand, we see that

$$\text{ord}_2\left(\frac{(x+1)\cdots(x+k)}{k!}\right) \leq \max_{1 \leq i \leq k} \text{ord}_2(x+i) \leq (\log(x+k))/\log 2. \tag{107}$$

Finally, we combine (104), (105), (106) and (107) to obtain (102).

As an application of Lemma 15, we prove

*Lemma 16.* If (14) holds, then

$$(x+1)\cdots(x+k) > \left(x + \frac{k}{2}\right)^k. \tag{108}$$

*Proof.* It is easy to check (108) for  $k \leq 4$ . Thus, we may assume that  $k \geq 5$ . Suppose that

$$(x+1)\cdots(x+k) \leq \left(x + \frac{k}{2}\right)^k. \tag{109}$$

By Lemma 11, we observe that  $A_1 - \binom{k}{1}\binom{k}{2} = \frac{k}{2}$ ,  $A_2 - \binom{k}{2}\binom{k}{2}^2 = \frac{k(k-1)(5k+2)}{24}$

and  $A_3 - \binom{k}{3}\binom{k}{2}^3 = \frac{k^2(k-1)(k-2)(2k+1)}{48}$  are positive. Further, we see from (103) and  $k \geq 5$  that  $x > k^2$ . Now, we derive from (109) that

$$\frac{1}{2}kx^{k-1} \leq \binom{k}{4}\binom{k}{2}^4 x^{k-4} \left(1 + \frac{1}{10} + \frac{1}{10^2} + \dots\right) \leq 5k^8 x^{k-4}/1728$$

which implies that  $172x^3 < k^7$ . Now, since  $x > k^2$ , we derive that  $k > 172$ . Then, we apply (103) to conclude that  $x > k^3$  and this is a contradiction.

For small values of  $y$  and  $k$ , the next two lemmas are useful. The first may be confirmed by direct checking.

*Lemma 17.*

$$P((x+1)\cdots(x+5)) \geq 31 \quad \text{for } 54 \leq x \leq 600$$

and

$$P((x + 1) \cdots (x + 6)) \geq 37 \quad \text{for } 120 \leq x \leq 600.$$

*Lemma 18.* Equation (14) implies that

$$y \geq \begin{cases} 17 & \text{if } k \in \{5, 6, 7, 8\} \\ 28 & \text{if } k = 9 \\ 32 & \text{if } k = 10 \end{cases}$$

*Proof.* By the arithmetic-geometric mean, we observe that

$$((x + 1) \cdots (x + k))^{1/k} < \frac{(x + 1) + \cdots + (x + k)}{k} = x + \frac{k + 1}{2}$$

which, together with (108), implies that

$$\lceil ((x + 1) \cdots (x + k))^{1/k} \rceil = x + \left\lceil \frac{k}{2} \right\rceil. \tag{110}$$

By applying again the arithmetic-geometric mean, we see from (87) and (92) that

$$\lceil ((y + 1) \cdots (y + 2k))^{1/k} \rceil \leq u + \left\lceil \frac{(k + 1)(2k + 1)}{3} \right\rceil. \tag{111}$$

Now, we combine (14), (110) and (111) to derive that

$$0 < x - u \leq \left\lceil \frac{(k + 1)(2k + 1)}{3} \right\rceil - \left\lceil \frac{k}{2} \right\rceil. \tag{112}$$

Let  $k \in \{5, 6, 7, 8\}$  and  $y \leq 16$ . Then, we see from (88) and (112) that  $x \leq 600$ . Now, we apply Lemma 17 to derive that  $x < 54$  if  $k \in \{5, 6, 7\}$  and  $x < 120$  if  $k = 8$  which, together with (103), imply that  $k = 5$  and  $x \geq 27$ . Therefore, we see from (88) and (112) that  $y(y + 11) < 54$  which implies that  $y \leq 3$ . Furthermore, we observe  $\lceil (y + 10)/4 \rceil - \lfloor y/4 \rfloor - \lfloor 10/4 \rfloor = 1$  if  $y = 2, 3$  to derive from Lemma 15 with  $\tau \geq 1$  that  $y \leq 1$ . Then, since the left hand side of (14) is not divisible by 31, we see that  $x \geq 31$ . Consequently, we observe that the left hand side of (14) is greater than the right hand side of (14).

Let  $k = 10$  and  $y \leq 31$ . Then, we observe from (88) and (112) that  $x \leq 1684$  which, by Lemma 15, implies that  $\tau = 0$ . On the other hand, we observe that  $\lceil (y + 20)/32 \rceil - \lfloor y/32 \rfloor - \lfloor 20/32 \rfloor = 1$  if  $y \geq 12$ . Consequently, we conclude that  $y \leq 11$  which implies that  $x \leq 424$  contradicting (103).

Let  $k = 9$  and  $y \leq 27$ . Then, by (88), (112) and (103), we observe that  $x \leq 1301$ ,  $x \geq 503$  and  $y \geq 14$ . Then, we observe  $\lceil (y + 18)/32 \rceil - \lfloor y/32 \rfloor - \lfloor 18/32 \rfloor = 1$  to apply Lemma 15 with  $\tau \geq 1$  to derive that  $x \geq 1015$  which implies that  $y > 22$ . If  $y \in \{23, 26, 27\}$ , we see  $\tau \geq 2$  to obtain from (102) that  $x \geq 2039$  which is not possible. If  $y = 24$ , then  $x \leq 1091$  and, by looking at the prime factors of 1086, 1091, 1093, we derive from (14) that  $x \leq 1076$  which is not possible, since the left hand side of (14) is less than the right hand side of (14). If  $y = 25$ , then  $x \leq 1159$  and we argue, as above by looking at the prime factors of 1151, 1159, 1162, to derive  $x \leq 1149$  for arriving at a contradiction. □

Now, we are ready to prove an inequality for (87) analogous to (108).

*Lemma 19.* Let  $u > 0$  if  $k = 3$  and

$$\theta = \begin{cases} 1 & \text{if } k \equiv 0 \pmod{3} \\ 2 & \text{if } k \equiv 1, 2 \pmod{3} \end{cases}$$

Then

$$\prod_{j=1}^k (u + j(2k - j + 1)) > \left( u + \frac{(k + 1)(2k + 1) - \theta}{3} \right)^k.$$

*Proof.* The Lemma can be verified for  $k \leq 4$ . By direct computations, we may assume that  $u \leq 21$  if  $k = 5$ ,  $u \leq 145$  if  $k = 6$  and  $u \leq 125$  if  $k = 7$ . Now, we refer to (88) and Lemma 18 to exclude these possibilities. Thus, we may suppose that  $k \geq 8$ . We set

$$h(\theta) = h(\theta, k) = ((k + 1)(2k + 1) - \theta)/3.$$

We observe that

$$h(\theta) \leq k(2k + 3)/3. \tag{113}$$

We suppose that

$$\prod_{j=1}^k (u + j(2k - j + 1)) \leq (u + h(\theta))^k. \tag{114}$$

By (113), we see that

$$\begin{aligned} (u + h(\theta))^k &\leq u^k + kh(\theta)u^{k-1} \\ &\quad + \binom{k}{2} \left( \frac{k(2k + 3)}{3} \right)^2 u^{k-2} + \dots + \binom{k}{k} \left( \frac{k(2k + 3)}{3} \right)^k. \end{aligned} \tag{115}$$

We put

$$D_2 = \frac{2}{45}k^5 - \frac{2}{9}k^4 - \frac{1}{6}k^3 + \frac{5}{18}k^2 + \frac{1}{15}k,$$

$$D_3 = \frac{k(k-1)(k-2)}{5670} (168k^5 + 32k^4 - 870k^3 - 1120k^2 - 423k - 72)$$

and

$$R' = \binom{k}{5} \left( \frac{k(2k + 3)}{3} \right)^5 u^{k-5} + \dots + \binom{k}{k} \left( \frac{k(2k + 3)}{3} \right)^k.$$

Further, we observe from (112), (103) and  $k \geq 8$  that  $u > k^2(2k + 3)/9$ . Therefore

$$R' \leq 2 \binom{k}{5} \left( \frac{k(2k + 3)}{3} \right)^5 u^{k-5}. \tag{116}$$

Thus, we derive from (114), (89), (115), (116) and Lemma 11 that

$$\begin{aligned} \left( \frac{\theta k}{3} \right) u^{k-1} &\leq D_2 u^{k-2} + D_3 u^{k-3} + \binom{k}{4} \left( \frac{k(2k + 3)}{3} \right)^4 u^{k-4} \\ &\quad + 2 \binom{k}{5} \left( \frac{k(2k + 3)}{3} \right)^5 u^{k-5}. \end{aligned} \tag{117}$$

We derive from (117) that  $u \leq 550$  if  $k = 8$ ,  $u \leq 1200$  if  $k = 9$  and  $u \leq 1500$  whenever  $k = 10$ . Now, we apply (88) and Lemma 18 to exclude these cases. Furthermore, the inequality (117) implies that  $u < 1950$  if  $k = 11$ ,  $u < 3980$  if  $k = 12$  and these cases are excluded by (112) and (103).

Thus, we may assume that  $k \geq 13$ . Then, we see from (112) and (103) that

$$u \geq \frac{7}{25} k^4. \tag{118}$$

Further, we observe that

$$D_2 < \frac{2}{45} k^5, \quad D_3 < \frac{171}{5670} k^8,$$

$$\binom{k}{4} \left( \frac{k(2k+3)}{3} \right)^4 \leq \frac{1}{24} \left( \frac{29}{39} \right)^4 k^{12} \leq \frac{13}{1000} k^{12}$$

and

$$2 \binom{k}{5} \left( \frac{k(2k+3)}{3} \right)^5 \leq \frac{1}{60} \left( \frac{29}{39} \right)^5 k^{15} \leq \frac{1}{250} k^{15}.$$

Therefore, we derive from (117) and  $\theta \geq 1$  that

$$u < \frac{2}{15} k^4 + \frac{513}{5670} \frac{k^7}{u} + \frac{39}{1000} \frac{k^{11}}{u^2} + \frac{3}{250} \frac{k^{14}}{u^3}$$

which, together with (118), implies that

$$\frac{7}{25} k^4 \leq u < \frac{2}{15} k^4 + k^3 + k^2.$$

Consequently, we conclude that  $k \leq 7$  which is a contradiction. □

As an immediate consequence of Lemma 19 and (110), we obtain the following improvement of (112).

**COROLLARY 3.**

Suppose that (14) is satisfied. Then, either  $x = 7$ ,  $y = 0$ ,  $k = 3$  or

$$x - u = f. \tag{119}$$

*Proof.* We may assume that  $u > 0$  if  $k = 3$ . Then, we see from Lemma 19 that

$$\left[ \left( \prod_{j=1}^k (u + j(2k - j + 1)) \right)^{1/k} \right] \geq u + [h(\theta)] = u + \left[ \frac{k(2k+3)}{3} \right]. \tag{120}$$

On the other hand, we see from the arithmetic-geometric mean and (92) that the left hand side of (120) is less than  $u + [(k+1)(2k+1)/3]$ . Consequently, we conclude that (120) holds with equality sign. Hence, the assertion (119) follows from (14), (87) and (110).

*Proof of Theorem 4.* We may assume that  $u > 0$  if  $k = 3$ . Then, we conclude from (14),

(87) and (119) that

$$(u + f + 1) \cdots (u + f + k) = \prod_{j=1}^k (u + j(2k - j + 1)). \tag{121}$$

We may verify, by direct computation, that (121) is not possible whenever  $k \leq 7$ . Thus, we may suppose that  $k \geq 8$ . Further, by (121), (93) and (89), we see that

$$(A_1(f) - B_1)u^{k-1} + (A_2(f) - B_2)u^{k-2} + \cdots + (A_k(f) - B_k) = 0.$$

Thus

$$(A_1(f) - B_1)u^{k-1} + (A_2(f) - B_2)u^{k-2} + (A_3(f) - B_3)u^{k-3} \leq (B_4 - A_4(f))u^{k-4} + R_3 \tag{122}$$

and

$$(B_1 - A_1(f))u^{k-1} + (B_2 - A_2(f))u^{k-2} + (B_3 - A_3(f))u^{k-3} \leq (A_4(f) - B_4)u^{k-4} + R_4 \tag{123}$$

where

$$R_3 = B_5 u^{k-5} + \cdots + B_k, \quad R_4 = A_5(f)u^{k-5} + \cdots + A_k(f). \tag{124}$$

By (119) and (103), we see that

$$u \geq \frac{22}{45} k(k+1)^2 \text{ if } k \geq 9 \quad \text{and} \quad u \geq \frac{1}{2}fk. \tag{125}$$

Thus, we derive from (124), (96), (98) and (125) that

$$R_3 \leq \frac{2}{145} k^5(k+1)^{10}u^{k-5} \quad \text{if } k \geq 9 \quad \text{and} \quad R_4 \leq \frac{17}{4} f^5 \binom{k}{5} u^{k-5}. \tag{126}$$

Next, we turn to the coefficient  $B_4 - A_4(f)$  in (122) and (123). By (91), we calculate  $A_4 = 22449$  if  $k = 8$  and  $A_4 = 157773$  if  $k = 10$ . Then, we derive from (94) and (99) that

$$(A_4(f) - B_4)/10^4 \leq \begin{cases} 30730 & \text{if } k = 8 \\ 487570 & \text{if } k = 10 \end{cases} \tag{127}$$

Similarly, by observing  $B_4 - A_4(f) \leq B_4 - \binom{k}{4}f^4 - \binom{k-1}{3}A_1f^3$ , we obtain

$$(B_4 - A_4(f))/10^5 \leq \begin{cases} 45990 & \text{if } k = 9 \\ 606420 & \text{if } k = 11 \end{cases} \tag{128}$$

and

$$(B_4 - A_4(f))/u^2 \leq \begin{cases} 11450 & \text{if } k = 12 \\ 7660 & \text{if } k = 13 \end{cases} \tag{129}$$

since, by (119) and (103),  $u \geq 3982$  if  $k = 12$  and  $u \geq 8060$  if  $k = 13$ . Next, we notice that

$$f \leq \frac{3}{4}k^2 \tag{130}$$

and

$$u \geq \frac{2}{3}k^4 \quad \text{if } k \geq 14. \tag{131}$$

First, we consider the case that  $k \equiv 2, 4 \pmod{6}$ . By Lemma 12 with  $\delta = -2/3$ , we



see that  $A_1(f) - B_1 = -k/2$  and

$$A_2(f) - B_2 < \frac{2}{45}k^5 - \frac{1}{3}k^4, \quad A_3(f) - B_3 < \frac{4}{135}k^8 - \frac{341}{2835}k^7,$$

since  $k \geq 8$ . Therefore, we derive from (123) that

$$\begin{aligned} \frac{1}{2}ku^{k-1} &< \left(\frac{2}{45}k^5 - \frac{1}{3}k^4\right)u^{k-2} + \left(\frac{4}{135}k^8 - \frac{341}{2835}k^7\right)u^{k-3} \\ &+ (A_4(f) - B_4)u^{k-4} + R_4. \end{aligned} \quad (132)$$

Then, we see from (132), (127) and (126) that  $u \leq 530$  if  $k = 8$  and  $u \leq 1300$  if  $k = 10$ . Now, we apply (88) and Lemma 18 to assume that  $k \geq 14$ . Further, we see from (96), (97), (130) and (126) that

$$A_4(f) \leq (1 + \phi_{14})f^4k^4/24 \leq k^{12}/25 \quad (133)$$

and

$$R_4 \leq 17f^5k^5u^{k-5}/480 \leq 9k^{15}u^{k-5}/1000. \quad (134)$$

Thus, we combine (132), (131), (133) and (134) to derive that

$$\frac{2}{5}k^4 \leq u < \frac{4}{45}k^4 + \frac{8}{135}\frac{k^7}{u} + \frac{2}{25}\frac{k^{11}}{u^2} + \frac{9}{500}\frac{k^{14}}{u^3} \leq \frac{4}{45}k^4 + k^3 + k^2$$

which implies that  $k \leq 4$ .

We argue, as above, to derive from (122) that

$$\frac{2}{3}ku^{k-1} \leq (B_4 - A_4(f))u^{k-4} + R_3 \quad (k \equiv 3 \pmod{6}), \quad (135)$$

$$\begin{aligned} \frac{1}{6}ku^{k-1} + \left(\frac{2}{45}k^5 + \frac{7}{72}k^4\right)u^{k-2} \\ \leq (B_4 - A_4(f))u^{k-4} + R_3 \quad (k \equiv 0 \pmod{6}) \end{aligned} \quad (136)$$

and

$$\left(\frac{2}{45}k^5 - \frac{7}{72}k^3\right)u^{k-2} \leq (B_4 - A_4(f))u^{k-4} + R_3 \quad (k \equiv 1, 5 \pmod{6}). \quad (137)$$

We see from (136), (137), (129) and (126) that  $u < 3982$  if  $k = 12$  and  $u < 8060$  if  $k = 13$ . On the other hand, we apply (119) and (103) to exclude these possibilities. Further, we see from (135), (137), (128) and (126) that  $u \leq 1300$  if  $k = 9$  and  $u \leq 3950$  if  $k = 11$ . In view of (88) and Lemma 18, the first case is not possible. Let  $k = 11$ . Then  $\tau = 0$  and  $35 \leq y \leq 52$ . Further, we observe  $[(y+22)/64] - [y/64] - [22/64] = 1$  if  $y \geq 42$  to derive from Lemma 15 that  $y \leq 41$ . Further, corresponding to every  $y$  with  $35 \leq y \leq 41$ , there is precisely one value of  $x$  given by (119) and (88). Finally, we count the power of 2 on both the sides of (14) in each of these seven cases to arrive at a contradiction.

Thus, we may assume that  $k \not\equiv 2, 4 \pmod{6}$  and  $k \geq 15$ . Then, we apply (131) to derive from (135), (136) and (137) that

$$\left(\frac{2}{45}k^5 - \frac{7}{72}k^3\right)u^{k-2} \leq (B_4 - A_4(f))u^{k-4} + R_3 \quad (k \not\equiv 2, 4 \pmod{6}). \quad (138)$$

Further, by (98) and (126), we see that

$$B_4 - A_4(f) < B_4 \leq 3k^{12}/40, \quad R_3 \leq 4k^{15}u^{t-5}/145. \quad (139)$$

Now, we combine (138), (139) and (131) to obtain

$$\frac{2}{45}k^5 \leq \frac{3}{40} \frac{k^{12}}{u^2} + \frac{4}{145} \frac{k^{15}}{u^3} + \frac{7}{72}k^3 \leq \frac{15}{32}k^4 + k^3$$

which implies that  $k \leq 12$ . □

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