

## On the ratio of values of a polynomial

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**Abstract.** For given positive integers  $a$  and  $b$ , the equation  $a(x+1)\dots(x+k) = b(y+1)\dots(y+k)$  in positive integers is considered. More general equations are also considered.

**Keyword.** Exponential diophantine equations.

### 1. Introduction

We shall always assume that  $a, b, x, y$  and  $k$  are positive integers satisfying  $b > a$  and  $x - y \geq k$ . We consider the equation

$$a(x+1)\dots(x+k) = b(y+1)\dots(y+k). \quad (1)$$

We re-write (1) as

$$u_k + A_1 u_{k-1} + \dots + A_k u_0 = 0 \quad (2)$$

where

$$u_n = ax^n - by^n, \quad 0 \leq n \leq k$$

and  $A_1, \dots, A_k$  are positive integers given by

$$(X+1)\dots(X+k) = X^k + A_1 X^{k-1} + \dots + A_k.$$

Let  $\alpha$  be a real number given by

$$ax^\alpha = by^\alpha.$$

Thus

$$\alpha = \left(\log \frac{b}{a}\right) \left(\log \frac{x}{y}\right)^{-1} > 0.$$

We prove:

**THEOREM 1.** *There exist effectively computable positive numbers  $C_1, C_2$  and  $C_3$  depending only on  $a$  and  $b$  such that (1) with  $k \geq C_1$  implies that*

$$C_2 k^2 x^{-1} < k - \alpha < C_3 \eta \quad (3)$$

where

$$\eta = \min(k^2 x^{-1}, x^{-1/3}).$$

In particular, (1) implies

$$0 < k - \alpha < 1, \quad k \geq C_1.$$

Thus we have:

COROLLARY 1. Equation (1) implies that either  $k < C_1$  or  $k = [\alpha + 1]$ .

For given positive integers  $a, b$  and  $k \geq 3$  such that the binary form  $aX^k - bY^k$  is irreducible over rationals, it follows from a theorem of Schinzel [3] that (1) has only finitely many solutions in  $x$  and  $y$ . We remark that this result is not effective.

We apply an estimate of Baker [1] on linear forms in logarithms to prove:

THEOREM 2. Equation (1) implies that  $\max(x, y, k)$  is bounded by an effectively computable number depending only on  $a, b$  and the greatest prime factor of  $xy$ .

It follows from theorem 1 that (1) implies

$$x \geq C_4 k^3$$

where  $C_4 > 0$  is an effectively computable number depending only on  $a$  and  $b$ . Further it follows from Crammer's conjecture on distance between consecutive primes that (1) implies

$$(\log x)^2 \geq C_5 k$$

where  $C_5 > 0$  is an absolute constant. We apply an inequality of van der Poorten [2] on  $p$ -adic linear forms in logarithms to obtain the following result.

THEOREM 3. Equation (1) implies that

$$\log x \leq C_6 k$$

where  $C_6 > 0$  is an effectively computable number depending only on  $a, b$  and the greatest prime factor of  $y(x - y)$ .

We state the following direct consequence of theorem 3.

COROLLARY 2. Equation (1) implies that  $\max(x, y)$  is bounded by an effectively computable number depending only on  $a, b, k$  and the greatest prime factor of  $y(x - y)$ .

Now we apply an argument of theorem 1 to a more general equation. For positive integers  $m$  and  $H$ , denote by  $S(m, H)$  the set of all polynomials

$$P(X) = B_0 X^m + B_1 X^{m-1} + \dots + B_m$$

where  $B_0 > 0, B_1 > 0$  and  $B_2, \dots, B_m$  are non-negative integers not exceeding  $H$ . For positive integers  $a_1, b_1, x_1$  and  $y_1$  with  $b_1 > a_1$  and  $x_1 > y_1$ , let  $\beta$  be a positive real number given by

$$a_1 x_1^\beta = b_1 y_1^\beta.$$

Then we have:

THEOREM 4. Let  $\delta > 0$  and  $P \in S(m, H)$ . Let  $a_1, b_1, x_1$  and  $y_1$  with  $b_1 > a_1$  and  $x_1 > y_1$  be positive integers. For  $l_1 \in \mathbb{Z}$  with

$$|l_1| \leq m^{-1} x_1^{m-1-\delta}.$$

suppose that

$$a_1 P(x_1) = b_1 P(y_1) + l_1. \quad (4)$$

There exist effectively computable positive numbers  $C_7, C_8$  and  $C_9$  depending only on  $\delta, a_1,$

$b_1$  and  $H$  such that for every  $m \geq C_7$ , we have

$$C_8 x_1^{-1} < m - \beta < \min(C_9 x_1^{-1}, 1). \tag{5}$$

Thus (4) implies

$$0 < m - \beta < 1, \quad m \geq C_7.$$

By taking  $P(X) = X^m + X^{m-1} + \dots + 1$ , we see from theorem 4 that

$$a_1 \frac{x_1^{m+1} - 1}{x_1 - 1} = b_1 \frac{y_1^{m+1} - 1}{y_1 - 1} + l_1$$

with  $m \geq C_7$  implies (5).

### 2. Linear forms in logarithms

We shall need the following results on linear forms in logarithms for the proof of theorems 2 and 3. Let  $\alpha_1, \dots, \alpha_n$  be non-zero rational numbers of heights not exceeding  $A_1, \dots, A_n$  respectively, where we assume that  $A_j \geq 3$  for  $1 \leq j \leq n$ . The height of a rational number  $E/F$  with  $(E, F) = 1$  is defined as  $\max(|E|, |F|)$ . Put

$$\Omega' = \prod_{j=1}^{n-1} \log A_j \quad \text{and} \quad \Omega = \Omega' \log A_n.$$

**THEOREM A** (Baker [1]). *There exist effectively computable absolute constants  $C_{10} > 0$  and  $C_{11} > 0$  such that the inequalities*

$$0 < |\alpha_1^{b_1} \dots \alpha_n^{b_n} - 1| < \exp(- (C_{10} n)^{C_{11} n} \Omega \log \Omega \log B)$$

have no solution in rational integers  $b_1, \dots, b_n$  of absolute values not exceeding  $B (\geq 3)$ .

**THEOREM B** (van der Poorten [2]). *Let  $p > 0$  be a prime number. There exist effectively computable absolute constants  $C_{12} > 0$  and  $C_{13} > 0$  such that the inequalities*

$$\infty > \text{ord}_p(\alpha_1^{b_1} \dots \alpha_n^{b_n} - 1) > (C_{12} n)^{C_{13} n} \Omega (\log B)^2 \frac{p}{\log p}$$

have no solution in rational integers  $b_1, \dots, b_n$  with absolute values at most  $B (\geq 3)$ .

### 3. Proof of theorem 1

Denote by  $c_1, c_2, \dots, c_{18}$  effectively computable positive numbers depending only on  $a$  and  $b$ . We may assume that  $k \geq c_1$  with  $c_1$  sufficiently large. Then, since  $x \neq y$ , observe that  $u_k u_{k-1} \neq 0$  (see lemma 1 of [4]). Further  $c_1 \leq k \leq x - y < x$  and (1) implies that none of  $(x+1), \dots, (x+k)$  is a prime number. Therefore it follows from the well-known results on difference between consecutive primes\* that

$$x \geq k^{3/2}. \tag{6}$$

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\* It is not necessary to use the results on difference between consecutive primes.

Now we sharpen inequality (6). Denote by  $d$  the greatest common divisor of  $(x+1) \dots (x+k)$  and  $(y+1) \dots (y+k)$ . By (1), we see that  $(x+1) \dots (x+k)/d$  divides  $b$ . In particular

$$x^k < (x+1) \dots (x+k) \leq bd.$$

By an argument of Erdos, the contribution in  $d$  from primes not exceeding  $2k$  is at most

$$k^k(x+k)^{\pi(2k)}$$

Further, by (1), the contribution in  $d$  of primes greater than  $2k$  does not exceed

$$(x-y+k)^{2k}((2k)!)^{-1}.$$

By (1), we see that  $ax^k < b(y+k)^k$  which implies

$$x-y \leq c_2 k^{-1}x + 2k. \tag{7}$$

Hence

$$x^k < bd \leq (c_3 k)^k x^{\pi(k)} \left( \max\left(\frac{x}{k^2}, 1\right) \right)^{2k}$$

which, together with (6), implies

$$x \geq c_4 k^3. \tag{8}$$

Combining (7) and (8), we have

$$x-y \leq 2c_2 k^{-1}x. \tag{9}$$

All the summands on the left side of (2) cannot be of the same sign. Therefore, since  $x > y$ , we see that  $u_k > 0$  which, together with (9), implies

$$x-y \geq c_5 k^{-1}x. \tag{10}$$

Suppose that  $u_{k-1} > 0$ . Observe that

$$u_k - xu_{k-1} = by^{k-1}(x-y).$$

Therefore, by (9) and (10), we see that

$$u_k \geq c_6 k^{-1}x^k.$$

Notice that

$$A_n \leq \binom{k}{n} k^n \leq k^{2n}, \quad 1 \leq n \leq k. \tag{11}$$

Re-write (1) as

$$u_k + A_1 u_{k-1} = -A_2 u_{k-2} - \dots - A_k u_0. \tag{12}$$

The absolute value of the left side of (12) is at least

$$\max(u_k, A_1 u_{k-1}) \geq c_6 k^{-1}x^k.$$

On the other hand, it follows from (11) and (8) that the absolute value of the right side of (12) is at most

$$c_7 k^4 x^{k-2}.$$

Comparing these estimates, we obtain

$$x \leq c_8 k^{5/2}.$$

Now we apply (8) to conclude that  $k \leq c_9$  which is not possible if  $c_1 > c_9$ .

Hence  $u_k > 0$  and  $u_{k-1} < 0$ . By continuity, we see that

$$0 < k - \alpha < 1. \tag{13}$$

Re-write (1) as

$$u_k = -A_1 u_{k-1} - A_2 u_{k-2} - \dots - A_k u_0.$$

Further

$$|u_{k-1}| = -u_{k-1} = b y^{k-1} (1 - (x/y)^{-(\alpha-k+1)}).$$

Now it follows from (9), (10), (13) and  $A_1 = \frac{k(k+1)}{2}$  that

$$c_{10}^{-1} (\alpha - k + 1) k x^{k-1} \leq |A_1 u_{k-1}| \leq c_{10} k x^{k-1}.$$

Similarly

$$|A_2 u_{k-2}| \leq c_{11} k^3 x^{k-2}$$

which, together with (8), implies that

$$|A_2 u_{k-2}| \leq c_{12} x^{k-1}.$$

Further, by (11) and (8),

$$|A_3 u_{k-3} + \dots + A_k u_0| \leq c_{13} k^6 x^{k-3} \leq c_{14} x^{k-1}.$$

Therefore  $u_k \leq 2c_{10} k x^{k-1}$ . Also, by (9), (10) and (13),

$$c_{15} \left( \frac{k-\alpha}{k} \right) x^k \leq u_k \leq c_{16} \left( \frac{k-\alpha}{k} \right) x^k.$$

Consequently  $k - \alpha < c_{17} k^2 x^{-1} < 1/4$ . Now notice that  $x^{-k} u_k \geq (2c_{10})^{-1} k x^{-1}$  and hence  $k - \alpha > c_{18} k^2 x^{-1}$ . Thus we have

$$c_{18} k^2 x^{-1} < k - \alpha < c_{17} k^2 x^{-1}$$

which, together with (8), implies (3). This completes the proof of theorem 1.

#### 4. Proof of theorems 2 and 3

Denote by  $v_1, \dots, v_7$  effectively computable positive numbers depending only on  $a, b$  and the greatest prime factor of  $xy$ . If  $k = 1$ , then we re-write (1) as

$$ax - by = b - a \neq 0$$

which, by theorem A, implies that  $\max(x, y) \leq v_1$ . Thus we may assume that  $k > 1$ . Further we apply again theorem A to obtain

$$x - y \geq x(\log x)^{-v_2}$$

which, together with (9), implies

$$k \leq (\log x)^{v_3}. \tag{14}$$

Let  $\delta$  be the least non-negative integer such that  $u_{k-\delta} \neq 0$ . Observe that  $\delta \in \{0, 1\}$ , since  $x \neq y$ . By theorem A,

$$|u_{k-\delta}| \geq x^{k-\delta} (k \log x)^{-v_4}$$

On the other hand, it follows from (2), (11) and (14) that

$$|u_{k-\delta}| \leq x^{k-\delta-1} (\log x)^{v_5}.$$

Comparing these estimates, we obtain

$$x \leq k^{v_6}$$

which, together with (14), implies that  $\max(x, y, k) \leq v_7$ . This completes the proof of theorem 2.

*Proof of theorem 3*

Denote by  $v_8, v_9, \dots$  effectively computable positive numbers depending only on  $a, b$  and the greatest prime factor of  $y(x - y)$ . We may assume that  $x \geq v_8$  with  $v_8$  sufficiently large, otherwise the assertion follows immediately. Suppose that

$$\log x > k.$$

It follows from (1) that  $x - y$  divides

$$\Delta = (b - a)(y + 1) \dots (y + k) \neq 0.$$

Thus, for a prime  $p$  dividing  $x - y$ , we have

$$\text{ord}_p(x - y) \leq \text{ord}_p(\Delta).$$

Let  $1 \leq n_0 \leq k$  satisfy

$$\text{ord}_p(y + n_0) \geq \text{ord}_p(y + n)$$

for  $n = 1, \dots, k$ . By theorem B, we have

$$\text{ord}_p(y + n_0) \leq v_9 (\log \log x)^3.$$

Therefore

$$\text{ord}_p(x - y) \leq \text{ord}_p(\Delta) \leq v_9 (\log \log x)^3 + v_{10} k.$$

Hence

$$\log(x - y) \leq v_{11} ((\log \log x)^3 + k).$$

On the other hand we see, by (10) which is also valid for  $k < c_1$ , that

$$\log(x - y) \geq \log x - \log k - v_{12}.$$

Comparing these estimates, we obtain  $\log x \leq v_{13} k$ . This completes the proof of theorem 3.

**5. Proof of theorem 4**

In the proof of theorem 4, we omit some details which already appear in the proof of theorem 1. Denote by  $c_{19}, c_{20}, \dots$  effectively computable positive numbers depending

only on  $\delta, a_1, b_1$  and  $H$ . We may assume that  $m \geq c_{19}$  with  $c_{19}$  sufficiently large. Let  $P \in S(m, H)$  be given by

$$P(X) = B_0 X^m + B_1 X^{m-1} + \dots + B_m.$$

Re-write (4) as

$$B_0 U_m + B_1 U_{m-1} + \dots + B_m U_0 = l_1 \tag{15}$$

where

$$U_n = a_1 x_1^n - b_1 y_1^n, \quad 0 \leq n \leq m.$$

Observe that (4) implies

$$c_{20} m^{-1} x_1 < x_1 - y_1 < c_{21} m^{-1} x_1. \tag{16}$$

Observe that

$$U_m - x_1 U_{m-1} = b_1 y_1^{m-1} (x_1 - y_1)$$

which implies

$$\max(|U_m|, |U_{m-1}|) \geq c_{22} m^{-1} x_1^{m-1}.$$

Thus all the summands on the left side of (15) cannot be of the same sign. Then, since  $x_1 > y_1$ , we see that  $U_m > 0$ . Suppose  $U_{m-1} > 0$ . Then

$$U_m \geq c_{23} x_1^{m-1}.$$

Further (15) implies that  $x_1 \leq c_{24}$  which, by (4), is not possible if  $c_{19}$  is sufficiently large.

Hence  $U_m > 0$  and  $U_{m-1} < 0$ . Thus

$$0 < m - \beta < 1. \tag{17}$$

Let  $N = \lceil (\log m)^2 \rceil$ . Using (16) and (17), we have

$$c_{25} \left( \frac{m-\beta}{m} \right) x_1^m < U_m < c_{26} \left( \frac{m-\beta}{m} \right) x_1^m,$$

$$c_{27} \left( \frac{\beta-m+1}{m} \right) x_1^{m-1} < |U_{m-1}| < c_{28} m^{-1} x_1^{m-1},$$

and

$$|U_{m-r}| \leq c_{29} r m^{-1} x_1^{m-r}, \quad 2 \leq r \leq N$$

$$|U_{m-r}| \leq c_{30} x_1^{m-r}, \quad N < r \leq m.$$

In view of these estimates, (15) implies (5). This completes the proof of theorem 4.

**Remark**

In theorem 4, the restriction on  $|l_1|$  can be relaxed to  $|l_1|$  does not exceed constant times  $m^{-1} x_1^{m-1}$ .

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